

## C50 Enumerative & Asymptotic Combinatorics

Stirling and Lagrange

Spring 2003

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This section of the notes contains proofs of Stirling's formula and the Lagrange Inversion Formula.

### Stirling's formula

**Theorem 1 (Stirling's Formula)**

$$n! \sim \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$$

**Proof** Consider the graph of the function  $y = \log x$  between  $x = 1$  and  $x = n$ , together with the piecewise linear functions shown in Figure 1.

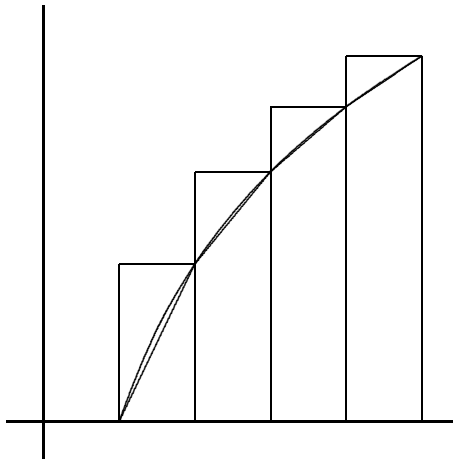


Figure 1: Stirling's formula

Let  $f(x) = \log x$ , let  $g(x)$  be the function whose value is  $\log m$  for  $m \leq x < m + 1$ , and let  $h(x)$  be the function defined by the polygon with vertices  $(m, \log m)$ , for  $1 \leq m \leq n$ . Clearly

$$\int_1^n g(x) dx = \log 2 + \cdots + \log n = \log n!.$$

The difference between the integrals of  $g$  and  $h$  is the sum of the areas of triangles with base 1 and total height  $\log n$ ; that is,  $\frac{1}{2} \log n$ .

Some calculus<sup>1</sup> shows that the difference between the integrals of  $f$  and  $g$  tends to a finite limit  $c$  as  $n \rightarrow \infty$ .

Finally, a simple integration shows that

$$\int_1^n f(x) dx = n \log n - n + 1.$$

We conclude that

$$\log n! = n \log n - n + \frac{1}{2} \log n + (1 - c) + o(1),$$

so that

$$n! \sim \frac{Cn^{n+1/2}}{e^n}.$$

To identify the constant  $C$ , we can proceed as follows. Consider the integral

$$I_n = \int_0^{\pi/2} \sin^n x dx.$$

Integration by parts shows that

$$I_n = \frac{n-1}{n} I_{n-2},$$

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<sup>1</sup>Let  $F(x) = f(x) - g(x)$ . The convexity of  $\log x$  shows that  $F(x) \geq 0$  for all  $x \in [m, m + 1]$ . For an upper bound we use the fact, a consequence of Taylor's Theorem, that

$$\log x \leq \log m + \frac{x-m}{m} \leq \log m + \frac{1}{m}$$

for  $x \in [m, m + 1]$ . Then

$$F(x) = \log x - \log m - \log \left(1 + \frac{1}{m}\right) (x-m) \leq \frac{1}{m} - \log \left(1 + \frac{1}{m}\right) \leq \frac{1}{2m^2},$$

where the last inequality comes from another application of Taylor's Theorem which yields  $\log(1+x) \geq x - x^2/2$  for  $x \in [0, 1]$ . Now  $\sum(1/m^2)$  converges, so the integral is bounded.

and hence

$$\begin{aligned} I_{2n} &= \frac{(2n)! \pi}{2^{2n+1} (n!)^2}, \\ I_{2n+1} &= \frac{2^{2n} (n!)^2}{(2n+1)!}. \end{aligned}$$

On the other hand,

$$I_{2n+2} \leq I_{2n+1} \leq I_{2n},$$

from which we get

$$\frac{(2n+1)\pi}{4(n+1)} \leq \frac{2^{4n} (n!)^4}{(2n)!(2n+1)!} \leq \frac{\pi}{2},$$

and so

$$\lim_{n \rightarrow \infty} \frac{2^{4n} (n!)^4}{(2n)!(2n+1)!} = \frac{\pi}{2}.$$

Putting  $n! \sim Cn^{n+1/2}/e^n$  in this result, we find that

$$\frac{C^2 e}{4} \lim_{n \rightarrow \infty} \left(1 + \frac{1}{2n}\right)^{-(2n+3/2)} = \frac{\pi}{2},$$

so that  $C = \sqrt{2\pi}$ .

## Lagrange Inversion

A formal power series over a field, with zero constant term and non-zero term in  $x$ , has an inverse with respect to composition. The associative, closure, and identity laws are obvious, and the rule for finding the inverse in characteristic zero is known as *Lagrange inversion*. We work over  $\mathbb{R}$  for convenience.

### The theorem

The basic fact can be stated as follows.

**Proposition 2** *Let  $f$  be a formal power series over  $\mathbb{R}$ , with  $f(0) = 0$  and  $f'(0) \neq 0$ . Then there is a unique formal power series  $g$  such that  $g(f(x)) = x$ ; the coefficient of  $y^n$  in  $g(y)$  is*

$$\left[ \frac{d^{n-1}}{dx^{n-1}} \left( \frac{x}{f(x)} \right)^n \right]_{x=0} / n!.$$

This can be expressed in a more convenient way for our purpose. Let

$$\phi(x) = \frac{x}{f(x)}.$$

Then the inverse function  $g$  is given by the functional equation

$$g(y) = y\phi(g(y)).$$

Then Lagrange inversion has the form

$$g(y) = \sum_{n \geq 1} \frac{b_n y^n}{n!},$$

where

$$b_n = \left[ \frac{d^{n-1}}{dx^{n-1}} \phi(x)^n \right]_{x=0}.$$

**Example: Cayley's Theorem** The exponential generating function for rooted trees satisfies the equation

$$T^*(x) = x \exp(T^*(x)).$$

With  $\phi(x) = \exp(x)$ , we find that the coefficient of  $y^n/n!$  in  $T^*(y)$  is

$$\left[ \frac{d^{n-1}}{dx^{n-1}} \exp(nx) \right]_{x=0} = n^{n-1}.$$

Now there are  $n$  ways to root a given tree; so the number of trees is  $n^{n-2}$ , proving Cayley's Theorem.

## Proof of the theorem

The proof of Lagrange's inversion formula involves a considerable detour. The treatment here follows the book by Goulden and Jackson. Throughout this section, we assume that the coefficients form a field of characteristic zero; for convenience, we assume that the coefficient ring is  $\mathbb{R}$ .

First, we extend the notion of formal power series to *formal Laurent series*, defined to be a series of the form

$$f(x) = \sum_{n \geq m} a_n x^n,$$

where  $m$  may be positive or negative. if the series is not identically zero, we may assume without loss of generality that  $a_m \neq 0$ , in which case  $m$  is the *valuation* of  $f$ , written

$$m = \text{val}(f).$$

We define addition, multiplication, composition, differentiation, etc., for formal Laurent series as for formal power series. In particular,  $f(g(x))$  is defined for any formal Laurent series  $f, g$  with  $\text{val}(g) > 0$ . (This is less trivial than the analogous result for formal power series. In particular, we need to know that  $g(x)^{-m}$  exists as a formal Laurent series for  $m > 0$ . It is enough to deal with the case  $m = 1$ , since certainly  $g(x)^m$  exists. If  $\text{val}(g) = r$ , then  $g(x) = x^r g_1(x)$ , and so  $g(x)^{-1} = x^{-r} g_1(x)^{-1}$ , and we have seen that  $g_1(x)^{-1}$  exists as a formal power series, since  $g_1(0)$  is invertible.

We denote the derivative of the formal Laurent series  $f(x)$  by  $f'(x)$ .

We also introduce the following notation:  $[x^n]f(x)$  denotes the coefficient of  $x^n$  in the formal power series (or formal Laurent series)  $f(x)$ . The case  $n = -1$  is especially important, as we learn from complex analysis. The value of  $[x^{-1}]f(x)$  is called the *residue* of  $f(x)$ , and is also written as  $\text{Res } f(x)$ .

Everything below hinges on the following simple observation, which is too trivial to need a proof.

**Proposition 3** *For any formal Laurent series  $f(x)$ , we have  $\text{Res } f'(x) = 0$ .*

Now the following result describes the residue of the composition of two formal Laurent series.

**Theorem 4 (Residue Composition Theorem)** *Let  $f(x)$ ,  $g(x)$  be formal Laurent series with  $\text{val}(g) = r > 0$ . Then*

$$\text{Res}(f(g(x))g'(x)) = r \text{Res}(f(x)).$$

**Proof** It is enough to consider the case where  $f(x) = x^n$ , since  $\text{Res}$  is a linear function.

Suppose that  $n \neq -1$ , so that the right-hand side is zero. Then

$$\text{Res}(g^n(x)g'(x)) = \frac{1}{n+1} \text{Res} \left( \frac{d}{dx} g^{n+1}(x) \right) = 0.$$

So consider the case where  $n = -1$ . Let  $g(x) = ax^r h(x)$ , where  $a \neq 0$  and  $h(0) = 1$ . Then

$$\begin{aligned} g'(x)g(x)^{-1} &= \frac{d}{dx} \log g(x) \\ &= \frac{d}{dx} (\log a + r \log x + \log h(x)) \\ &= \frac{r}{x} + \frac{d}{dx} \log h(x), \end{aligned}$$

so

$$\text{Res } g'(x)g(x)^{-1} = r = r \text{Res } x^{-1}.$$

Note that we have cheated slightly in the first line of this argument:  $\log g(x)$  may not exist as a formal Laurent series; but it is the case that for if the equation  $f(x) = g(x)h(x)$  holds for formal Laurent series, then

$$f'(x)/f(x) = g'(x)/g(x) + h'(x)/h(x),$$

and in the preceding argument it is the case that  $\log h(x)$  exists (this is obtained by substituting  $y = h(x) - 1$  in  $\log(1 + y)$ ) and its derivative is  $h'(x)/h(x)$ . Consider this point carefully; an error here would lead to the incorrect conclusion that  $\text{Res}(g'(x)/g(x)) = 0$ .

From the Residue Composition Theorem, we can prove a more general version of Lagrange Inversion.

**Theorem 5 (Lagrange Inversion)** *Let  $\phi$  be a formal power series with  $\text{val}(\phi) = 1$ , Then the equation*

$$g(x) = x\phi(g(x))$$

*has a unique formal power solution  $g(x)$ . Moreover, for any Laurent series  $f$ , we have*

$$[x^n]f(g(x)) = \begin{cases} \frac{1}{n}[x^{n-1}](f'(x)\phi(x)^n) & \text{if } n \geq \text{val}(f) \text{ and } n \neq 0, \\ f(0) + \text{Res}(f'(x)\log(\phi(0)^{-1}\phi(x))) & \text{if } n = 0. \end{cases}$$

**Proof** Let  $\Phi(x) = x/\phi(x)$ , so that  $\Phi(g(x)) = x$  and  $\text{val}(\Phi(x)) = 1$ . Then  $g$  is the inverse function of  $\Phi$ .

We have

$$\begin{aligned} [x^n]f(g(x)) &= \text{Res}x^{-n-1}f(g(x)) \\ &= \text{Res}\Phi(y)^{-n-1}\Phi'(y)f(y), \end{aligned}$$

where we have made the substitution  $x = \Phi(y)$  (so that  $y = g(x)$ ) and used the Residue Composition Theorem.

For  $n \neq 0$ , we have

$$\begin{aligned} [x^n]f(g(x)) &= -\frac{1}{n}[y^{-1}]f(y) (\Phi(y)^{-n})' \\ &= \frac{1}{n}[y^{-1}]f'(y)\Phi(y)^{-n} \\ &= \frac{1}{n}[y^{n-1}]f'(y)\phi(y)^n. \end{aligned}$$

Here, in the second line, we have used the fact that

$$\text{Res}(f'(x)g(x)) = -\text{Res}(f(x)g'(x)),$$

a consequence of the fact that  $\text{Res}(f(x)g(x))' = 0$ ; in the third line we use the fact that  $\Phi(x) = x/\phi(x)$ .

For  $n = 0$ , we have

$$\begin{aligned} [x^0]f(g(x)) &= [y^0]f(y) - [y^{-1}]f(y)\phi'(y)\phi(y)^{-1} \\ &= f(0) + \text{Res}(f'(y)\log(\phi(y)\phi^{-1}(0))), \end{aligned}$$

using the same principle as before and the fact that

$$(\log(\phi(y)\phi^{-1}(0)))' = \phi'(y)\phi(y)^{-1}.$$

Taking  $f(x) = x$  in this result gives the form of Lagrange Inversion quoted earlier.

We proceed to an application, also taken from Goulden and Jackson, of the Residue Composition Theorem.

**Example: a binomial identity** We use the Residue Composition Theorem to prove that

$$\sum_{k=0}^n \binom{2n+1}{2k+1} \binom{j+k}{2n} = \binom{2j}{2n}.$$

We begin with the sum of the odd terms in  $(1+x)^{2n+1}$ :

$$\sum_{k=0}^n \binom{2n+1}{2k+1} x^{2k} = \frac{1}{2x} ((1+x)^{2n+1} - (1-x)^{2n+1}).$$

Call the right-hand side of this equation  $f(x)$ . Now, if  $S$  is the sum that we want to evaluate, then

$$\begin{aligned} S &= [y^{2n}](1+y)^j \sum_{k=0}^n \binom{2n+1}{2k+1} (1+y)^k \\ &= \text{Res} y^{-(2n+1)} (1+y)^j f((1+y)^{1/2}). \end{aligned}$$

Now we do the following rather strange thing: make the substitution  $y = z^2(z^2 - 2)$ . Then  $\text{val}(y(z)) = 2$ , and  $(1+y)^{1/2} = 1 - z^2$ . So the Residue Composition Theorem gives

$$\begin{aligned} S &= \text{Res}(z^2 - 1)^{2j} \left( \frac{1}{(z^2 - 2)^{2n+1}} - \frac{1}{z^{4n+2}} \right) z \\ &= \text{Res}(z^2 - 1)^{2j} z^{-(4n+1)} \\ &= [z^{4n}](z^2 - 1)^{2j} \\ &= \binom{2j}{2n}, \end{aligned}$$

as required. (In the second line we have used the fact that  $(z^2 - 2)^{-(2n+1)}$  is a formal power series and so its residue is zero.)