

Pym Tautological Talk

outline:

- 1) A_g and \bar{A}_g
- 2) Tautological ring
- 3) Torelli cycle
- 4) Pym ~~tautological~~ cycle
- 5) GRR
- 6) Algorithm on \bar{H}_g .

(Joint w/ Yaav Len and Sam Molcho.)

1) A_g and \bar{A}_g

- $A = V/\Lambda$ abelian variety of $\dim = g$.
[Eg. $A = \text{Jac}(C)$.]

- Principal polarisation: ~~homological~~ homological equivalence class of well-behaved ample divisor.

(Rigidities. Also \exists non-algebraic \mathbb{C} -tori.)

[Eg. Θ on $\text{Jac}(C)$. Needed to recover C .]

- A_g = moduli space of PPAV of $\dim g$.

- Construction: $A = \mathbb{C}^g/\Lambda$ with:

$$\Lambda = \mathbb{Z}^g + \boxed{\mathbb{Z}} \cdot \mathbb{Z}^g \quad \text{~~matrix~~ } \mathbb{Z}\text{-column span of } \mathbb{Z}$$

with $Z \in \text{Mat}_{g \times g}(\mathbb{C})$:

- Z symmetric
- $\text{Im } Z$ positive definite.

$H_g = \text{Space of } Z = \text{"Siegel upper half space"}$

$$A_g = H_g / \text{SP}(2g, \mathbb{Z})$$

- $\dim A_g = \dim H_g = \boxed{\binom{g+1}{2}}$.

A_g smooth DM stack, non-proper.

Alternative construction via Hilbert Schemes:
 $\Psi_L^3: A_c \hookrightarrow \mathbb{P}^{3g-1}$

- Compactification: there are many.
- Hiding details/technicalities, there is a moduli space

A^{trop}
 A_g

[Ash-Mumford-Rapoport-Tai '75]
 [Faltings-Chai '90]

with no preferred core complex structure:

choice $(\Sigma \rightarrow A_g^{trop}) \rightsquigarrow (\overline{A}_g^\Sigma \supseteq A_g)$ toroidal compactification

- What does \overline{A}_g^Σ parametrize? Fiddly.
- Semiabelian varieties (plus extra data):

*1

$$0 \rightarrow T \rightarrow G \rightarrow X \rightarrow 0$$

\parallel
 G_m^r

\hookrightarrow "degeneration data"
 \uparrow abelian, $\dim = g-r$.
 [Néron '61, Grothendieck '72]

eg: If $C \rightsquigarrow C_0$ by acquiring r nodes, then $\text{Jac}(C) \rightarrow \text{Jac}(C_0)$:

$$0 \rightarrow G_m^r \rightarrow \text{Jac}(C_0) \rightarrow \text{Jac}(\tilde{C}_0) \rightarrow 0$$

- Every \overline{A}_g^Σ carries a universal semiabelian variety:

(*)2

$$\begin{array}{ccc} \Sigma & \xrightarrow{q} & \bar{\Sigma} \\ U_g & \xleftarrow{e} & A_g \end{array}$$

(Stable under base change to a refinement.)

- Some carry a compactified universal family.

Idea: Compactify (*)1 to ~~the universal family~~ on expansion.

~~the universal family~~

■ Mumford degeneration: polyhedral decomposition of real tons.



Can reverse logic and use this to construct Σ , by combinatorial SS reduction.

- [strata of \bar{A}_g^1 explicit: tons bundle over abelian variety bundle over A_{g-1} .]

2) Tautological ring

- Look back at (*)2. Define:

$$\begin{array}{l} E = e^* \Omega_g \quad (\text{Hodge bundle, } \text{rk} = g.) \\ \lambda_i = c_i(E) \quad \blacksquare \end{array}$$

[Torelli:  $Mg \xrightarrow{\text{Tor}} Ag$  $\text{Tor}^k E = E$]
 $C \mapsto \text{Jac}(C)$

- Hodge bundle preserved under pullback along $\bar{A}_g^{\Sigma'} \rightarrow \bar{A}_g^{\Sigma}$

• Defⁿ: The tautological ring

$R^*(\bar{A}_g) \subseteq CH^*(\bar{A}_g)$.

is the subring generated by $\lambda_1, \dots, \lambda_g$.
 Doesn't depend on Σ (dropped from notation).

Relatively simple, e.g. no boundary classes.
 In fact: [Van der Geer '96]

• Th^m [vdG]: AS a graded ring:

$R^*(\bar{A}_g) = \mathbb{Q}[\lambda_1, \dots, \lambda_g] / ((1 + \lambda_1 + \dots + \lambda_g)(1 - \lambda_1 + \dots + (-1)^g \lambda_g))$

(Equivalent to $ch_{2k}(E) = 0$ for $k \geq 1$.)

Pf: GRR for universal \mathbb{Q} divisor to find relation.

Use $\lambda_1 \dots \lambda_g \neq 0$ and Properties of Gorenstein rings to show this is only relation. \square

$\Rightarrow R^*(\bar{A}_g)$ based by Suorefree monomials in x_i .

$\square R^*(\bar{A}_g)$ Gorenstein: restricting of intersection pairing is perfect.

[Not automatic: $R^*(\bar{M}_{g,n})$ is rarely Gorenstein, cf. Canning]

• Produces tautological projection:

$$\gamma \in \text{CH}^k(\bar{A}_g) \rightsquigarrow \langle \gamma, - \rangle: R^{(g+1)-k}(\bar{A}_g) \rightarrow \mathbb{Q}$$

$$\rightsquigarrow \text{taut}(\gamma) \in R^k(\bar{A}_g).$$

• Goal: Compute tautological projection of natural cycles on \bar{A}_g .

[Our techniques also apply to the tautological projection of \bar{A}_g , cf. Canning-Molcho-Oprea-Pandharipande]

3) Torelli cycle

• Torelli \square map extends:

$$\bar{M}_g \xrightarrow[\text{Tor}]{} \bar{A}_g^\Sigma, \quad \text{Tor}^* x_i = x_i$$

(Sometimes need to blowup \bar{M}_g ; it's a question of whether tropical Torelli respects the cone complex structures. \square But not necessary for $\Sigma = 1\text{st}$ or 2nd Voronoi, for what it's worth.)

• Computing $\text{tauA}(\text{Tor}^*[\overline{M}_g])$



Computing $\int_{\overline{M}_g} \lambda_{i_1} \dots \lambda_{i_k}$ (squarefree monomial)

• Faber (1997): algorithm (and code) for this.

Uses Mumford's GRR formula for λ_i .

Boundary classes appear \Rightarrow recursive.

* Key trick: Forget markings at end of each recursion step.

Valid since λ_i pulled back from $\overline{M}_{g,0}$.

Crucial to avoid explosion of combinatorial complexity. (Remember this!).

4) Prym cycle

• Jacobians nicest PPAVS.
Pryms second-nicest.

• $f: C \rightarrow B$ étale double cover of smooth curves.

$$f^* \text{Jac}(B) \subseteq \text{Jac}(C)$$

[Mumford '74]

~~□~~ $\text{Prym}(f) \subseteq \text{Jac}(C)$

complementary
abelian subvariety.

- $0 \rightarrow \text{Pym}(f) \rightarrow \text{Jac}(C) \xrightarrow{N_{M_f}} \text{Jac}(B) \rightarrow 0$ (up to extension by $\mathbb{Z}/2\mathbb{Z}$)
- $f^* \text{Jac}(B) \times \text{Pym}(f) \rightarrow \text{Jac}(C)$
 $\begin{matrix} +1 \text{ e.space} & -1 \text{ e.space} \\ \text{for } \mathbb{Z} & \text{for } \mathbb{Z} \end{matrix}$

- $\dim \text{Pym}(f) = \dim \text{Jac}(C) - \dim \text{Jac}(B)$
 $= g_C - g_B$
 $= g_B - 1$ (RH $\Rightarrow g_C = 2g_B - 1$)

[Why double covers? Why étale?
 Don't really need, but the polarisation isn't principal.
 our work applies, but formula less nice.]

- $\boxed{\text{Hg}}$ = Hurwitz space of étale double covers $C \rightarrow B, g_B = g$.

[$t: \text{Hg} \rightarrow \text{Mg}$ finite, degree = $(2^{2g} - 1)/2$.]

- $\left. \begin{matrix} \text{Hg} \xrightarrow{\text{Pym}} \text{Ag}_{-1} \\ (C \xrightarrow{f} B) \mapsto \text{pym}(f) \end{matrix} \right\}$

By tropical geometry, lifts to compactification:

$$\overline{\text{Hg}}^{\mathbb{Z}} \xrightarrow{\text{Pym}} \overline{\text{Ag}}_{-1}^{\mathbb{Z}}$$

[unlike Torelli, actually have to blowup $\overline{\text{Hg}}$ in all known cases.]

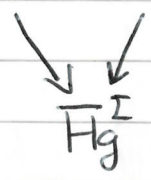
• computing $\text{tau}_t(\text{Prym}^*[Hg])$



Computing $\int_{Hg} \text{Prym}^* \lambda_{i_1} \dots \text{Prym}^* \lambda_{i_k}$

These are Prym lambda classes; consider

$$C \xrightarrow{f} B$$



There are three associated abelian varieties, each with their own Hodge bundle:

Jac(C)	E_C
Jac(B)	E_B
$\text{Prym}(f)$	\tilde{E}

They are related: 

$$0 \rightarrow \tilde{E} \rightarrow E_C \rightarrow E_B \rightarrow 0$$

and:

$\text{Prym}^* \lambda_i = C_i(\tilde{E})$

- All well-defined on \bar{H}_g , so can discard Σ and work there. (Projection formula.)
- Idea: Use projection formula for $t: \bar{H}_g \rightarrow \bar{M}_g$.

$$\hat{E} = E_C - E_B = E_C - t^*E$$
- Hope: Express E_C in terms of E_B . Then everything difficult will be pulled back via t .

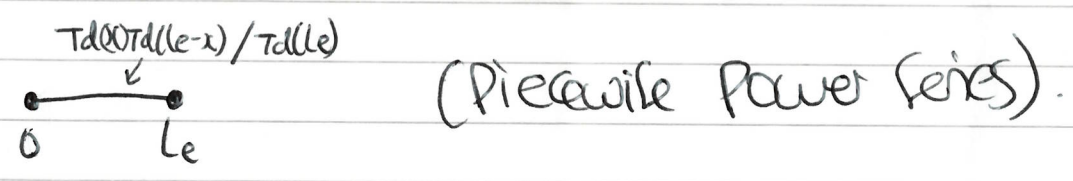
5) GRR

- $$\begin{array}{ccc}
 C & \xrightarrow{f} & B \\
 \rho_C \downarrow & & \downarrow \rho_B \\
 & \bar{H}_g &
 \end{array}$$

$$\begin{array}{l}
 E_C = \rho_{C*} \omega_{P_C} \\
 E_B = \rho_{B*} \omega_{P_B}
 \end{array}
 \left. \vphantom{\begin{array}{l} E_C \\ E_B \end{array}} \right\} \begin{array}{l} \text{Compute each} \\ \text{using GRR} \\ \text{and compare} \end{array}$$

$f \text{ log etale} \Rightarrow f^* \omega_B = \omega_C$

Key difference is Todd class: $Td(T_{P_C})$ vs. $Td(T_{P_B})$.



$\Rightarrow Td(T_{P_C}) \neq Td(T_{P_B})$ only along locus of nodes where f is ramified.

Key Simple

For etale double covers this locus is simple:

- Ramified \Rightarrow mult-2
- NO ramification \Rightarrow node cannot be cut markings separating (RH parity)

\Rightarrow only one boundary divisor contributes:

$$\Delta_0 = \left[\begin{array}{c} \circlearrowleft^2 \\ \downarrow \\ \circlearrowleft \\ \text{g-1} \end{array} \right],$$

$$L: \overline{Hg}_{-1, (2)}^2 \rightarrow \overline{Hg}$$

$$L_x(1) = 2\Delta_0$$

Allows us to compare E_C and $E_B = t^*E$.

Th^m: In \overline{Hg} we have:

$$ch(\tilde{E}^v) = t^* ch(E^v) - 1 +$$

$$\frac{1}{2} \int L_x \left(\sum_{k \geq 1} \frac{(2^{2k} - 1) B_{2k}}{(2k)!} (\psi_1^{2k-2} - \psi_1^{2k-3} \psi_2 + \dots + \psi_2^{2k-2}) \right)$$

For M_2 , equivalent to chiodo's formula.
 But ours applies to arbitrary G-covers.
 Already different from chiodo for $G=M_2$.

Essentially, poly appearing in Mumford's GRR calculation.
 Universal polynomial for Todd classes of codim. 2 loci

6) Implementation

- Remains to integrate Chern classes of \tilde{E} .
- Same idea: projection formula for t .

• Problem 1: $\Delta_0 \neq t^* S_0$.

In fact:

$$t^* S_0 = 2 \left[\begin{array}{c} \text{Diagram 1} \\ \downarrow \\ \text{Diagram 2} \end{array} \right] + \left[\begin{array}{c} \text{Diagram 3} \\ \downarrow \\ \text{Diagram 4} \end{array} \right] + \left[\begin{array}{c} \text{Diagram 5} \\ \downarrow \\ \text{Diagram 6} \end{array} \right]$$

(A₀)

(can prove using extended piecewise polynomials)

• So we have to work directly on $\bar{H}g$.

We can do this, but the combinatorics is complicated. which brings us to:

• Problem 2: There is no forgetful map:

$$\bar{H}g_{-1, (2)2} \rightarrow \bar{H}g_{-1}$$

Thus, unlike Faber, we can't control the number of markings in the recursion. The combinatorics explodes. ($g=6$ is bare cake!)

Still, hope to find a way to	one working example	perhaps I can	return on this in a future talk
---	--------------------------------	--------------------------	--

• ~~*~~ Solution: Instead of multiplying out RHS directly, use invariance of lambda classes.

$$\left. \begin{aligned} L^* \text{Ch}^v(\tilde{E}) &= \text{Ch}^v(\tilde{E}) && \text{(loop gluing)} \\ \mathbb{P}_1^* \text{Ch}^v(\tilde{E}) &= \text{Ch}^v(\tilde{E}_1) + \text{Ch}^v(\tilde{E}_2) && \text{(bridge gluing)} \end{aligned} \right\}$$

This vastly improves the efficiency.

- The algorithm is computer-implemented and effective for $g \leq 10$.

(obstacle for $g \geq 11$ is actually enumerating the basis of $R^*(A_{g-1})$.)

- In $g=6$ (first nontrivial case) we need Donagi's result:

$$\text{Prym}_*[\bar{F}_6] = 27 \cdot [\bar{A}_5]$$

- In $g=7$ we ~~find~~ find:

$$\text{twt Prym}_*[\bar{F}_7] = \frac{\boxed{691}}{178,541,140,377,600} (10 \lambda_1 \lambda_2 + \lambda_3)$$

- In $g=8$ we find (PTO):

total Pym_g [F₁₈] =

$$- \frac{691}{158,703,235,891,200} \lambda^7$$

$$+ \frac{69,791}{856,997,473,812,480} \lambda_1 \lambda_6$$

$$- \frac{131,981}{856,997,473,812,480} \lambda_2 \lambda_5$$

$$- \frac{21,421}{612,141,052,723,200} \lambda_3 \lambda_4$$

$$+ \frac{11,747}{194,772,153,139,200} \lambda_1 \lambda_2 \lambda_4$$

• Etc.

$g=9$: degree 12, 10 terms.

$g=10$: degree 18, 21 terms.