

MAS400: Solutions 4

Throughout this sheet F will denote an arbitrary field unless otherwise stated.

1. Let $f_1 = xy - 1$ and $f_2 = y^2 - 1$ be elements of $\mathbb{Q}[x, y]$. Show that $\{f_1, f_2\}$ is *not* a Gröbner basis for $\langle f_1, f_2 \rangle$, with respect to \leq_{lex} and $x > y$.

If $\{f_1, f_2\}$ is a Gröbner basis then $f \text{ rem } (f_1, f_2) = f \text{ rem } (f_2, f_1)$ for all $f \in \mathbb{Q}[x, y]$. However, in lectures we saw that $x^2y + xy^2 + y^2 \text{ rem } (f_1, f_2) = x + y + 1$, while $x^2y + xy^2 + y^2 \text{ rem } (f_2, f_1) = 2x + 1$ (the latter calculation was left as an exercise).

Alternatively: Newer technology (Theorem T) indicates we must consider whether $S(f_1, f_2) \text{ rem } (f_1, f_2) = 0$. We have $\text{lcm}(\text{lm } f_1, \text{lm } f_2) = xy^2$, so that $S(f_1, f_2) = \frac{xy^2}{xy}f_1 - \frac{xy^2}{y^2}f_2 = yf_1 - xf_2 = -y + x$. Multivariate division of $S(f_1, f_2)$ by (f_1, f_2) or (f_2, f_1) gives $q_1 = q_2 = 0$ and remainder $x - y \neq 0$, regardless of the monomial order chosen, so that $\{f_1, f_2\}$ is not a Gröbner basis for $\langle f_1, f_2 \rangle$.

(Note that in this case it is fairly easy to spot a suitable linear combination $f = p_1f_1 + p_2f_2$ of f_1 and f_2 such that $f \text{ rem } (f_1, f_2) \neq 0$; Theorem T and the S -polynomials just provides a systematic way of finding suitable linear combinations.)

2. Let $G = \{g_1, \dots, g_t\} \subseteq F[x_1, \dots, x_n]$ with $0 \notin G$. Suppose that, with respect to a fixed monomial order, we have:

$$\forall f \in \langle g_1, \dots, g_t \rangle, \quad f \text{ rem } (g_1, \dots, g_t) = 0.$$

Show that G is a Gröbner basis for $\langle G \rangle$.

We consider the algorithm **MultivariateDivision**. On each iteration of the while-loop we may add $\text{lt } p$ to r , or leave r unchanged. But each iteration of the while-loop also strictly decreases $\text{mdeg } p$, so no term that is added to r can cancel previous terms. For $f \in \langle G \rangle$ the eventual remainder is zero, and so the else-part of the if-loop in the while-loop is never executed; therefore at all times during the execution of **MultivariateDivision** we have $p = 0$ or that $\text{lt } p$ is divisible by (at least) one of the $\text{lt } g_i$. If $f \neq 0$ then upon starting the first iteration of the while loop we have $p = f$, so that $\text{lt } f$ is divisible by (at least) one of the $\text{lt } g_i$, and thus $\text{lt } f \in \langle \text{lt } g_1, \dots, \text{lt } g_t \rangle$. Therefore $\langle \text{lt}(\langle G \rangle) \rangle \leq \langle \text{lt } g_1, \dots, \text{lt } g_t \rangle$, and so G is a Gröbner basis for $\langle G \rangle$.

3. Let R be a ring (commutative with 1). Show that each ideal of R is finitely generated if and only if there are no infinite ascending chains of ideals:

$$I_0 < I_1 < I_2 < I_3 < \cdots.$$

NB: I have proved a very similar result in lectures.

IF I is not finitely generated then we define $I_0 = 0$, and I_m for $m \geq 1$ recursively as follows: We pick $r_m \in I \setminus I_{m-1}$. Since $I_{m-1} = \langle r_1, \dots, r_{m-1} \rangle$ is finitely generated and I is not, we can always pick such an r_m for all m . Then $I_0 < I_1 < I_2 < \cdots$ is a strictly ascending chain of R -ideals contained in I .

If each ideal is finitely generated we let $I_0 \leq I_1 \leq I_2 \leq \cdots$ be an ascending chain of ideals (possibly with equalities at various points). Now $I = \bigcup_{n=0}^{\infty} I_n$ is also an ideal (exercise), and so I is finitely generated, say $I = \langle r_1, \dots, r_s \rangle$. For $i \in \{1, \dots, s\}$ there exists $n_i \in \mathbb{N}$ such that $r_i \in I_{n_i}$. Let $N = \max(n_1, \dots, n_s)$. For $n \geq N$ we have $I_n \leq I$ (since I is the union of the I_n) and $I \leq I_n$ (since I_n contains a generating set for I); thus $I_n = I$. Therefore $I_N = I_{N+1} = I_{N+2} = \cdots$, and no strictly infinite ascending chain of ideals exists.